

Learning Unordered Tree Contraction Patterns in Polynomial Time

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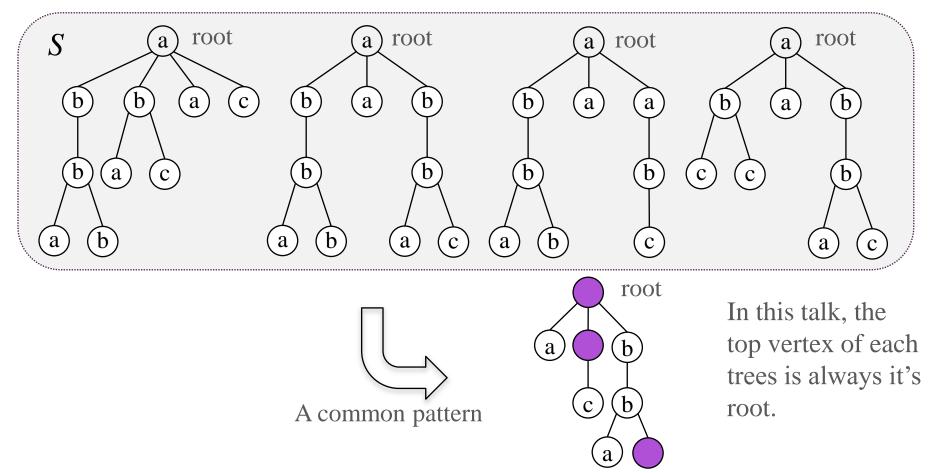
- 1. Backgrounds & motivations
- 2. Preliminaries
 - Tree contraction pattern (TC-pattern)
- 3. Time complexity of the TC-pattern matching problem
- 4. The minimal language problem for TC-patterns
- 5. Conclusions and future work

+ Backgrounds & motivations

- Increase of tree-structured data
 - Web documents
 - XML files etc..
- Discovery common characteristic tree-structured patterns from treestructured database
- Application
 - Classification of a tree-structured data set

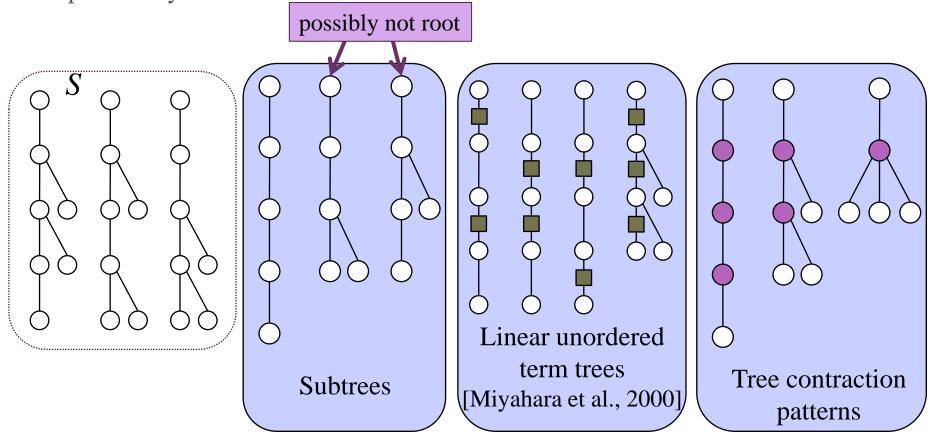
+ Backgrounds & motivations

• A graph pattern expression common to given tree structured data



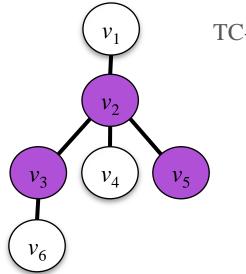
+ Backgrounds & motivations

A difference between tree contraction patterns and term trees which previously studied.



• A <u>tree contraction pattern</u> (<u>TC-pattern</u>) is a triplet $t = (V_p E_p U_t)$ where

- V_t is a vertex set,
- E_t is an edge set, and
- U_t is a subset of V_t , whose elements are called contractible vertices. Below, purple vertices indicate contractible vertices.
- $(V_p E_t)$ is a tree with a specified root $r_t \in V_t$.



TC-pattern *t*

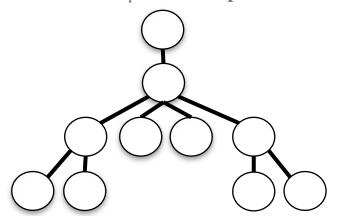
$$V_{t} = \{v_{1}, v_{2}, v_{3}, v_{4}, v_{5}, v_{6}\}$$

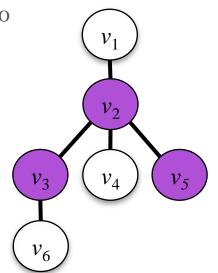
$$E_{t} = \{\{v_{1}, v_{2}\}, \{v_{2}, v_{3}\}, \{v_{2}, v_{4}\}, \{v_{2}, v_{5}\}, \{v_{3}, v_{6}\}\}$$

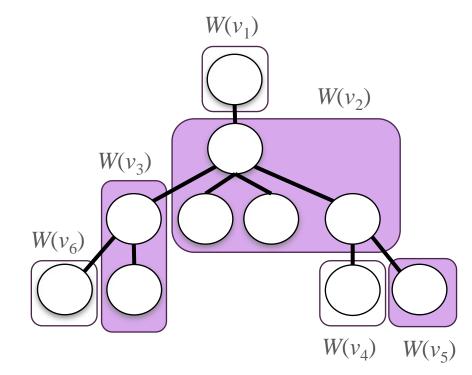
$$U_{t} = \{v_{2}, v_{3}, v_{5}\}$$

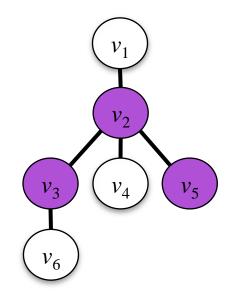
$$r_{t} = v_{1}$$

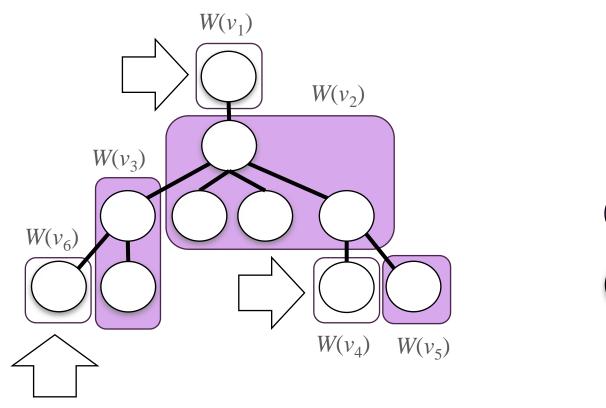
- A tree $T = (V_T, E_T)$ with root $r_T \underline{matches}$ a TC-pattern t with root r_t , if there is a partition of V_T , $\{W(v_1), \dots, W(v_m)\}$ for $v_1, \dots, v_m \in V_t$, such that
 - 1. for $v \in V_t \setminus U_t$, W(v) includes exactly one vertex,
 - 2. for any $v \in V_t$, any pair of W(v) is connectied,
 - 3. $W(r_t)$ includes r_T , and
 - 4. the tree obtained from *T* by merging all vertices in W(v) into one vertex for each $v \in U_t$ is isomorphic to *T*.



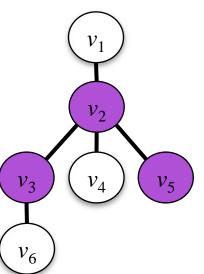


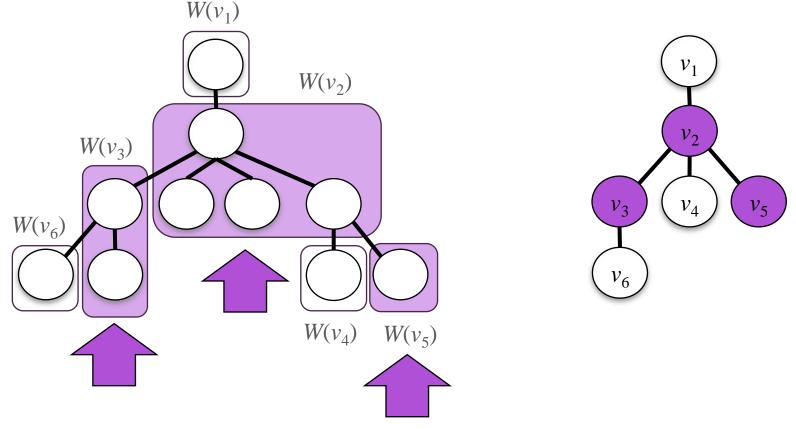




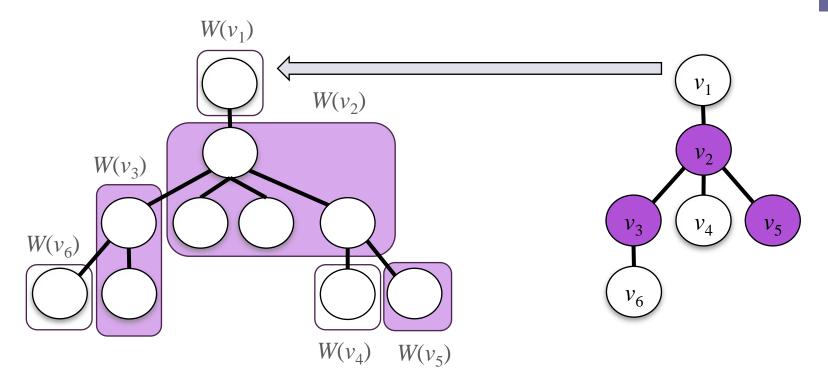


1. For $v \in V_t \setminus U_t$, W(v) includes exactly one vertex.

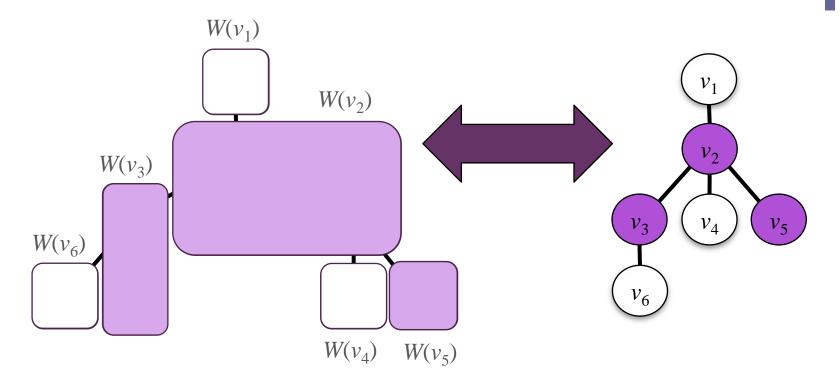




2. For any $v \in V_t$, the subtree induced by W(v) is connected.



3. $W(r_t)$ includes r_T .



4. The tree obtained from T by merging all vertices in W(v) into one vertex for each $v \in U_t$ is isomorphic to *t*.

Time complexity of the TC-pattern matching problem

TC-pattern matching problem

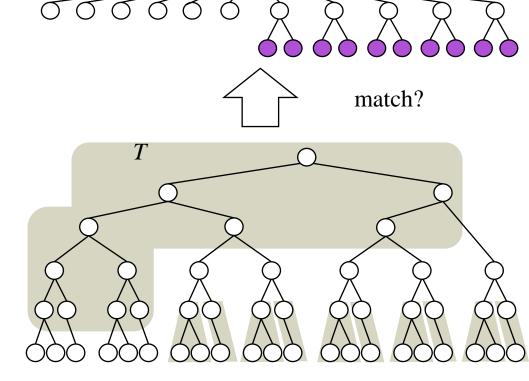
Input: a rooted unordered tree *T*, a TC-pattern *t*.

Question: *T* matches *t* ?

Theorem

TC-pattern matching problem is NP-complete.

Proof: Transform from X3C.



In this paper, we consider <u>a subclass of TC-patterns whose matching</u> problem can be solved in polynomial time.

Time complexity of the TC-pattern matching problem

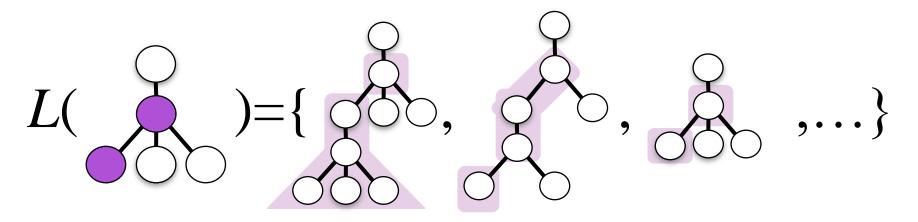
Theorem

We assume that the degree of every contractible vertex in TC-patterns is bounded by a constant *d*. Then TC-pattern matching problem for a given TC-pattern *t* and a given tree *T* is solved in $O(nN^{\max\{d-1,1.5\}})$ time, where $n=|V_t|$ and $N=|V_T|$.

Method: Dynamic Programing.

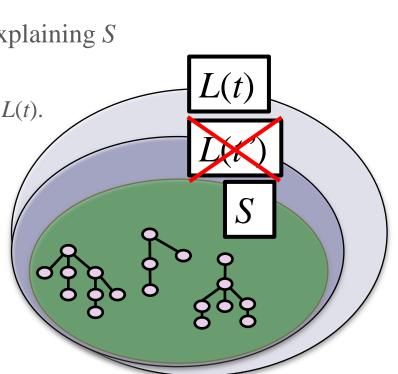
For every vertex v of T, we compute a unique label (a collection of subsets of t) by using the labels of the children of v.

- The TC-pattern language L(t)
 - A representation power of TC-pattern *t*.



• The TC-pattern language L(t) is defined as the set of all trees which match t.

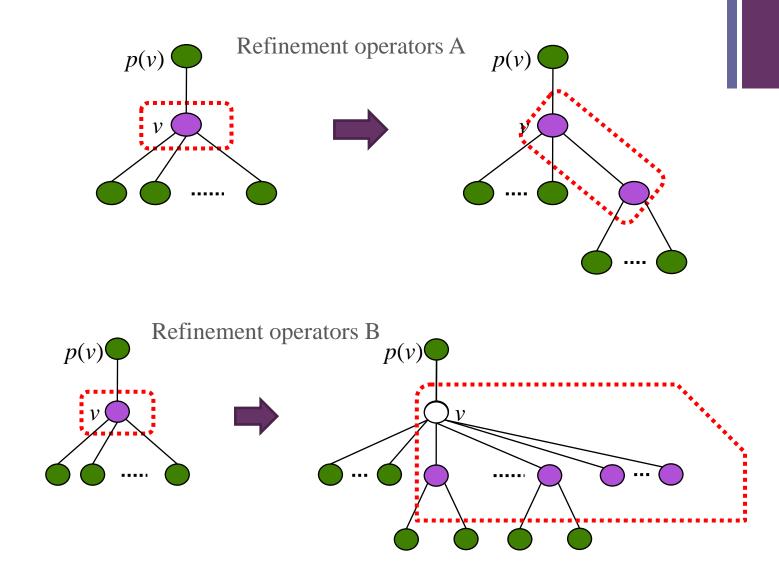
- MINimal Language (MINL) problem for TC-patterns
 - Instance: A set of rooted unordered trees $S = \{T_1, T_2, ..., T_m\}$
 - Problem: Find a minimally generalized TC-pattern explaining *S*.
- Def. <u>A minimally generalized TC-pattern t</u> explaining S
 - 1. L(t) contains all trees in S.
 - 2. There is no TC-pattern *t*' such that $S \subseteq L(t') \subsetneq L(t)$.
- From a data mining point of view, this problem is a problem for searching a given dataset for only one specialized common pattern.

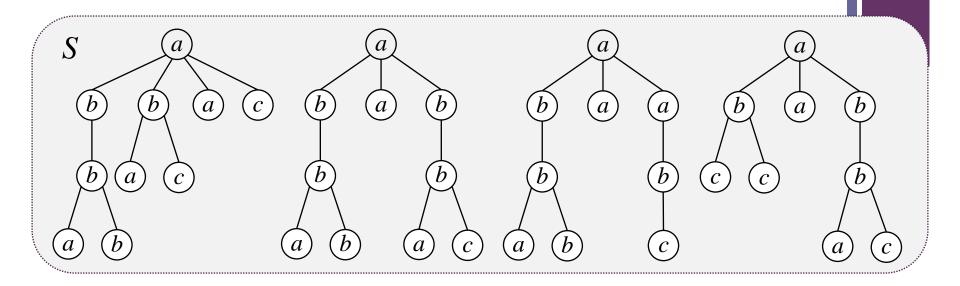


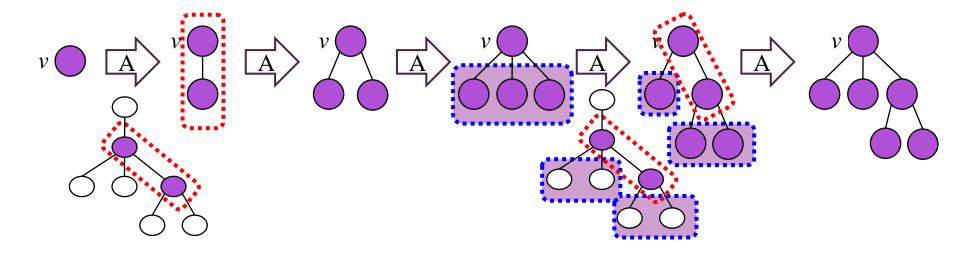
An idea to compute the MINL problem

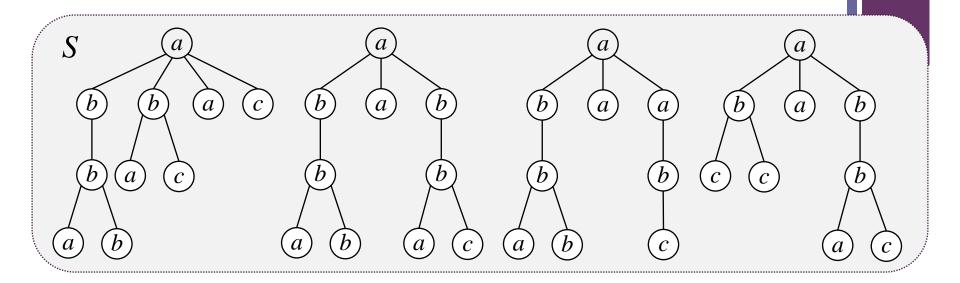
- Starting from the most generalized TC-pattern.
 - The most generalized TC-pattern is a TC-pattern consisting of only one contractible vertex.
- Trying to specialize TC-pattern t to provide a more specialized TC-pattern t' which satisfies the next conditions.
 - $S \subseteq L(t') \subseteq L(t)$, and
 - if there is no such t', output t.

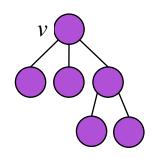
Next, we show two refinement operators which are used in this refinement process.

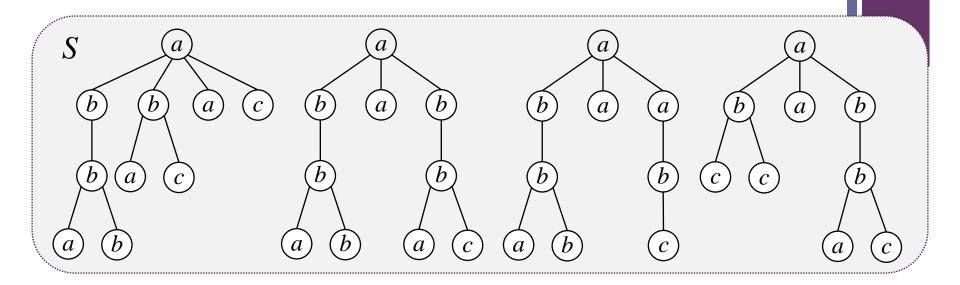


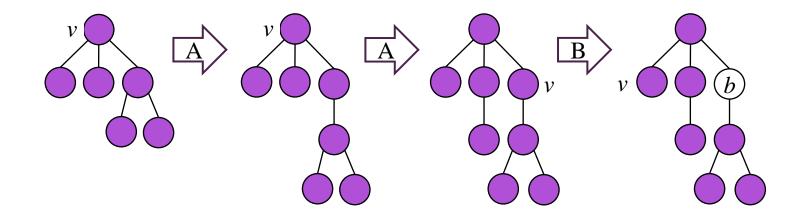


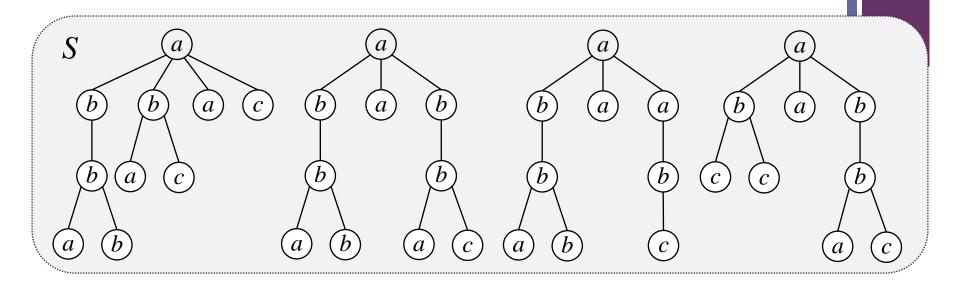


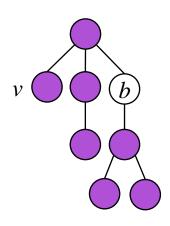


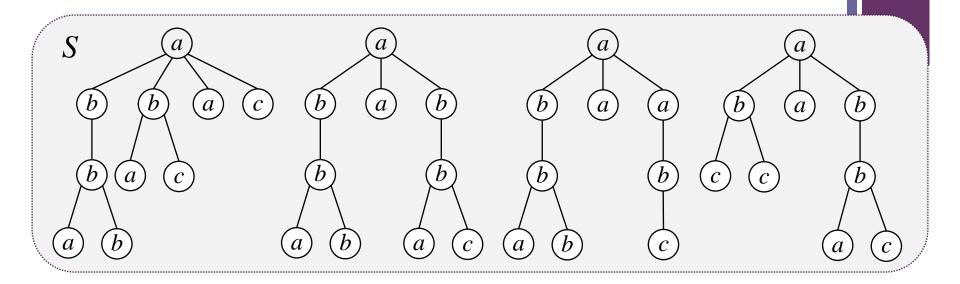


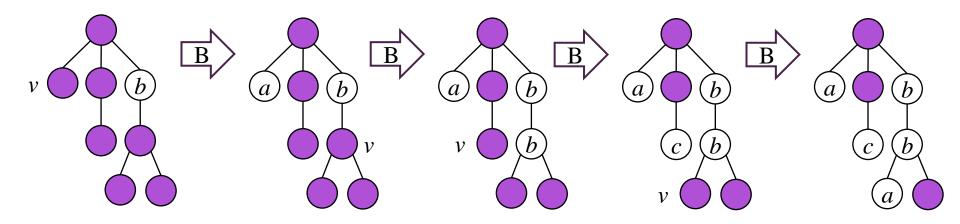












Theorem

- We assume that there are infinitely many vertex labels in \sum , and that the degree of every contractible vertex in TCpatterns is bounded by constant *d*.
- Minimal language problem for TC-patterns for a given set of trees *S* is computed in

 $O(nN_{\min}^{d+1}N_{\max}^{\max\{d-1,1.5\}}|S||\sum(S)|) \text{ time,}$ where $N_{\min} = \min_{T \in S} |V_T|$, $N_{\max} = \max_{T \in S} |V_T|$, and $\sum(S) = \{\delta \in \sum | \delta \in S \}$.

+ Conclusions

- A learning model on computational learning theory: A polynomial time inductive inference from a positive data
- Theorem[Angluin, '80, Shinohara, '82]: If a class *C* has finite thickness, and the membership and the minimal language (MINL) problems for *C* are solvable in polynomial time, class *C* is polynomial time inductively inferable from positive data.

Corollary

The class of TC-patterns such that <u>the degree of every contractible</u> <u>vertex in it is bounded by a constant *d* is polynomial time inductively inferable from positive data.</u>

+ Future work

- Experiment of our algorithm
 - Web document
 - Sugar chain data etc..
- Development of more fast algorithm to find a minimally generalized TC-pattern.
- To consider <u>graph contraction</u> patterns (GC-patterns) based on <u>tree</u> <u>contraction</u> patterns (TC-patterns).
 - outerplanar graph
 - bounded treewidth graph etc..